



Computer Architecture 10

Redundant Number Systems



Redundant Number Systems

- ▶ Conventional radix- r systems use $[0,r-1]$ digit set
 - radix-10 → 0,1,2,3,4,5,6,7,8,9
- ▶ If the digit set (in radix- r system) contains more than r digits, the system is redundant
 - radix-2 → 0,1,2 or -1,0,1
 - radix-10 → 0,1,2,3,4,5,6,7,8,9,10,11,12,13
 - radix-10 → 0,1,2,3,4,5,6,7,8,9,ڭ,ڭ,ڭ
 - radix-10 → -6,-5,-4,-3,-2,-1,0,1,2,3,4,5
- ▶ Conversions between Redundant Number Systems is a simple digit-serial process



Redundancy

- ▶ Redundancy – representation of numbers is not unique
- ▶ Redundancy may result from narrowing the range of represented values (e.g. 1's compl.), but number interpretation may be complex
- ▶ Redundancy may result from adopting the digit set wider than radix and the number interpretation is conventional
- ▶ Physical representation of redundant numbers may be not trivial.



Conversions

- ▶ Conversion is a carry-propagate (slow) process

11	9	17	10	12	18	radix-10, digit set [0,18]	
11	9	17	10	12	18	$18 = 10+8$	
11	9	17	10	13	8	$13 = 10+3$	
11	9	17	11	3	8	$11 = 10+1$	
11	9	18	1	3	8	$18 = 10+8$	
11	10	8	1	3	8	$10 = 10+0$	
12	0	8	1	3	8	$12 = 10+2$	
1	2	0	8	1	3	8	radix-10, digit set [0,9]

11	9	17	10	12	18	radix-10, digit set [0,18]	
11	9	17	10	12	18	$18 = 20-2$	
11	9	17	10	14	-2	$14 = 10+4$	
11	9	17	11	4	-2	$11 = 10+1$	
11	9	18	1	4	-2	$18 = 20-2$	
11	11	-2	1	4	-2	$11 = 10+1$	
12	1	-2	1	4	-2	$12 = 10+2$	
1	2	1	-2	1	4	-2	radix-10, digit set [-6,5]



Decomposition

- ▶ Redundant numbers with $[0,m]$ digit set can be represented by two numbers of $[0,n]$ digit sets, where $m=2n$
- ▶ Conversion requires ordinary addition of two such numbers with $[0,n]$ digit set representation

11 9 17 10 12 18	radix-10, digit set [0,18]
<hr/>	
9 9 9 9 9 9	
+	
2 0 8 1 3 9	
<hr/>	radix-10, digit set [0,9]
1 2 0 8 1 3 8	

- ▶ Decomposed representation is, of course, not unique, but the sum amounts to correct result

Radix-2 [0,2] digit set numbers

- ▶ Redundant binary numbers may be coded with bit-fields, e.g:
 - 0: (0,0),
 - 1: (0,1) or (1,0),
 - 2: (1,1)
- ▶ Decomposed form is convenient and efficient representation of redundant binary numbers

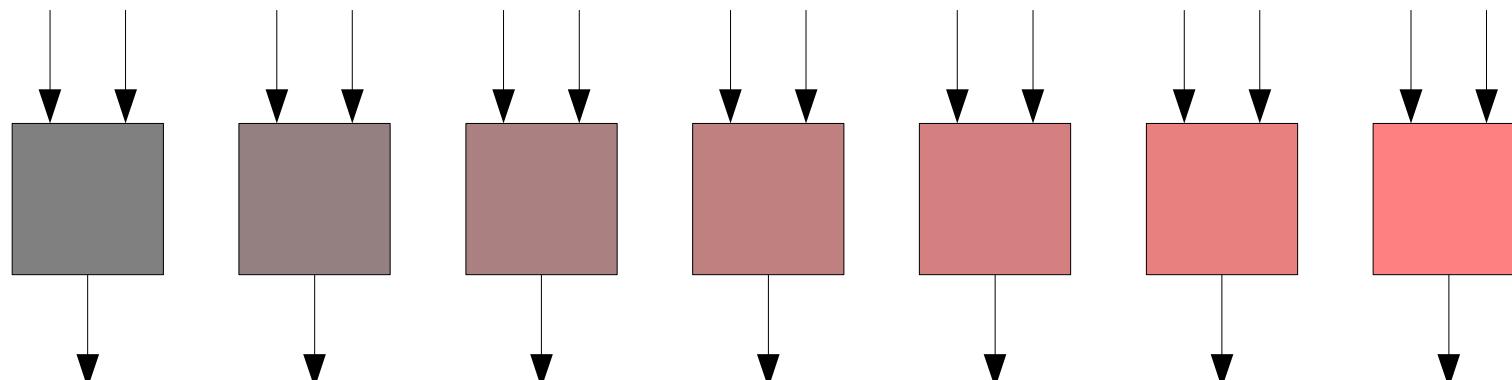
1	1	2	0	2	0		radix-2, digit set [0,2]
+ 0	0	1	0	1	0		
1 0 0 0 1 0 0							radix-2, digit set [0,1]



Carry-free addition

- ▶ Carry-free → no carry propagation, all digit additions can be done simultaneously
- ▶ Carry-free addition is possible with widening of the digit set

$$\begin{array}{rcccccc} 1 & 2 & 3 & 4 & 5 & 6 \\ 4 & 5 & 6 & 7 & 8 & 9 \\ \hline 5 & 7 & 9 & 11 & 13 & 15 \end{array} \quad \begin{array}{l} \text{radix-10, digit set [0,9]} \\ \text{radix-10, digit set [0,9]} \\ \text{radix-10, digit set [0,18]} \end{array}$$



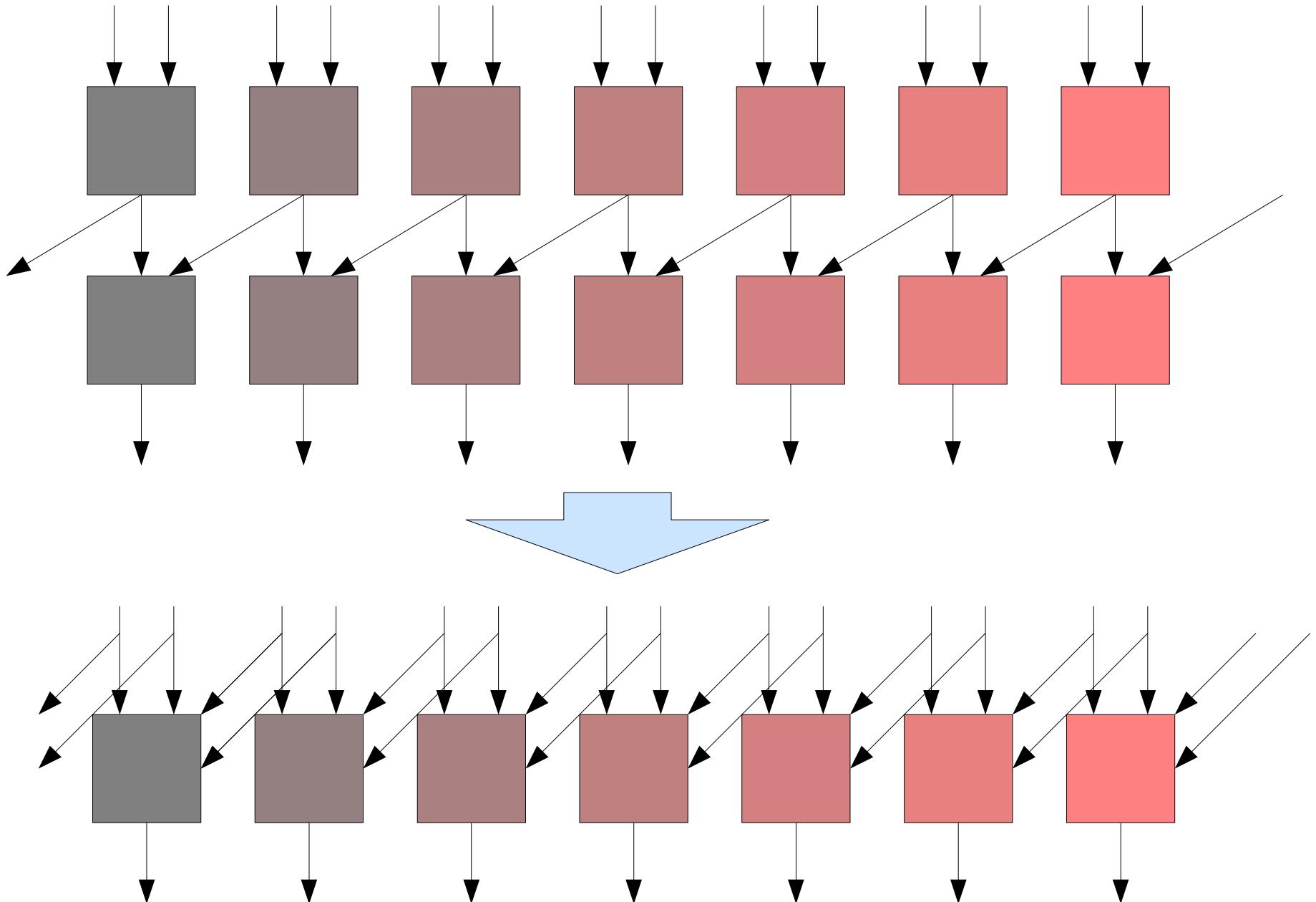


Two-stage carry-free addition

- ▶ Reduction of digit set by carry propagation by only one position

11 9 17 10 12 18	radix-10, digit set [0,18]
6 12 9 10 8 18	radix-10, digit set [0,18]
17 21 26 20 20 36	radix-10, digit set [0,36]
↓ ↓ ↓ ↓ ↓ ↓	
7 11 16 0 10 16	Intermediate sums [0,16]
↖ ↖ ↖ ↖ ↖ ↖	
1 1 2 1 2	Transfer digits [0,2]
8 12 18 1 12 16	Sum [0,18]

Two-stage carry-free addition

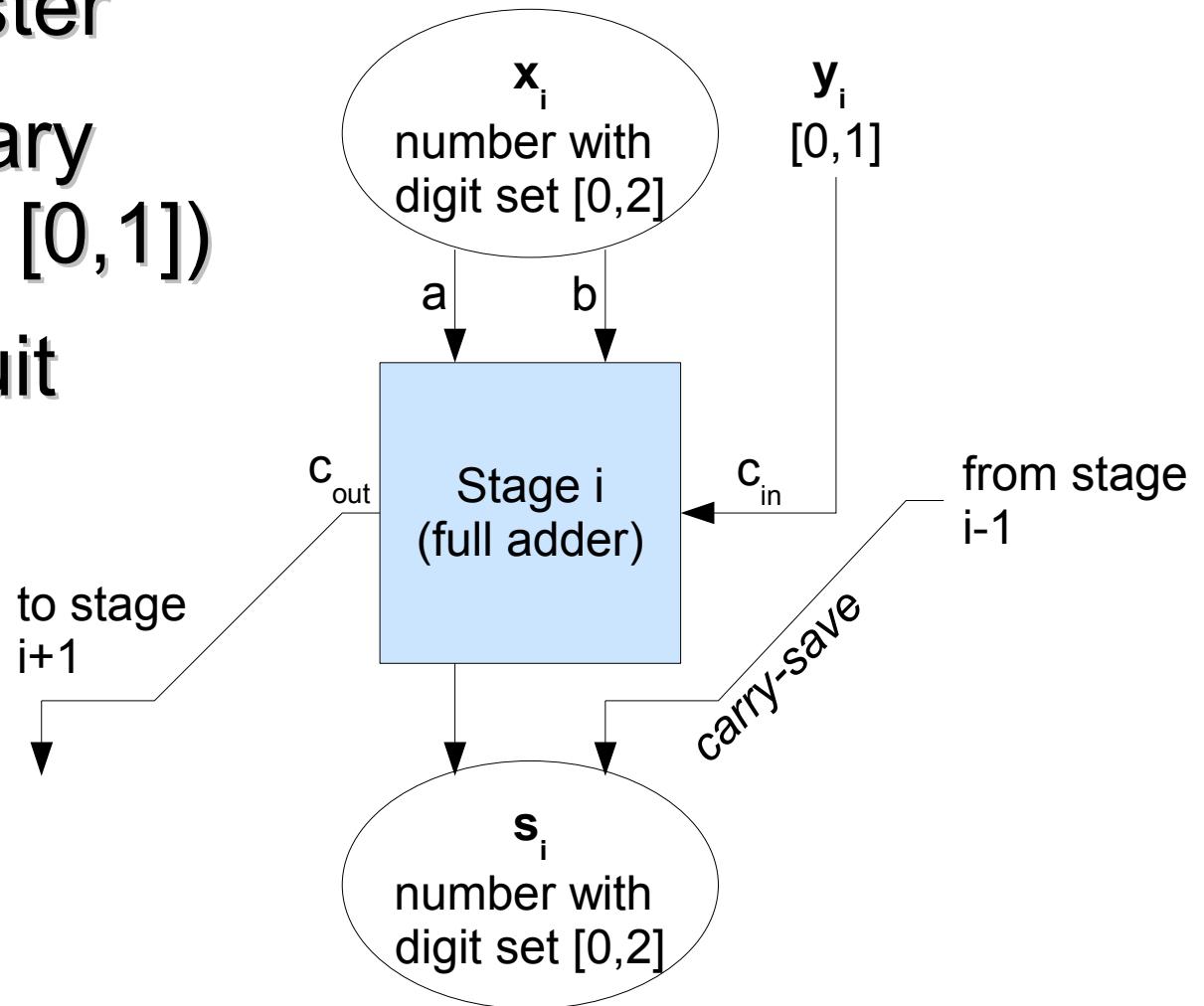


Redundant binary additions

$+ \begin{array}{ccccccc} 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 & 1 & 0 \end{array}$	1 st number, radix-2, digit set [0,1]
$= \begin{array}{ccccccc} 0 & 1 & 2 & 1 & 1 & 1 \end{array}$	2 nd number, radix-2, digit set [0,1]
$+ \begin{array}{ccccccc} 0 & 1 & 1 & 1 & 0 & 1 \\ 0 & 2 & 3 & 2 & 1 & 2 \end{array}$	Temp. sum, radix-2, digit set [0,2]
$= \begin{array}{ccccccc} 0 & 0 & 1 & 0 & 1 & 0 \end{array}$	3 rd number, radix-2, digit set [0,1] ([0,2])
$= \begin{array}{ccccccc} 0 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 2 & 3 & 2 & 1 & 2 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ 0 & 0 & 1 & 0 & 1 & 0 \end{array}$	Temp. sum, radix-2, digit set [0,3]
$= \begin{array}{ccccccc} 0 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \end{array}$	Temp. sums [0,1]
$+ \begin{array}{ccccccc} 0 & 1 & 1 & 2 & 0 & 2 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{array}$	Transfer digits [0,1]
$= \begin{array}{ccccccc} 1 & 1 & 3 & 0 & 3 & 1 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ 1 & 1 & 1 & 0 & 1 & 1 \end{array}$	Temp. sum [0,2]
$+ \begin{array}{ccccccc} 0 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 3 & 0 & 3 & 1 \end{array}$	4 th number, radix-2, digit set [0,1] ([0,2])
$= \begin{array}{ccccccc} 1 & 2 & 1 & 1 & 1 & 1 \end{array}$	Temp. sum, radix-2, digit set [0,3]
$= \begin{array}{ccccccc} 0 & 0 & 1 & 0 & 1 & 0 \\ 1 & 2 & 1 & 1 & 1 & 1 \end{array}$	Temp. sums [0,1]
$+ \begin{array}{ccccccc} 1 & 0 & 0 & 1 & 1 & 1 & 1 \end{array}$	Transfer digits [0,1]
$= \begin{array}{ccccccc} 1 & 0 & 0 & 1 & 1 & 1 & 1 \end{array}$	Result of multiple addition [0,2]
$\leftarrow \leftarrow \leftarrow \leftarrow \leftarrow \leftarrow$	Carry-propagate conversion
$\boxed{1 \ 0 \ 0 \ 1 \ 1 \ 1 \ 1}$	Result of multiple addition [0,1]

Carry-save adder

- ▶ Carry is saved in next stage register
- ▶ Circuit adds 3 binary numbers (digit set [0,1])
- ▶ 3/2 reduction circuit





Signed-digit numbers

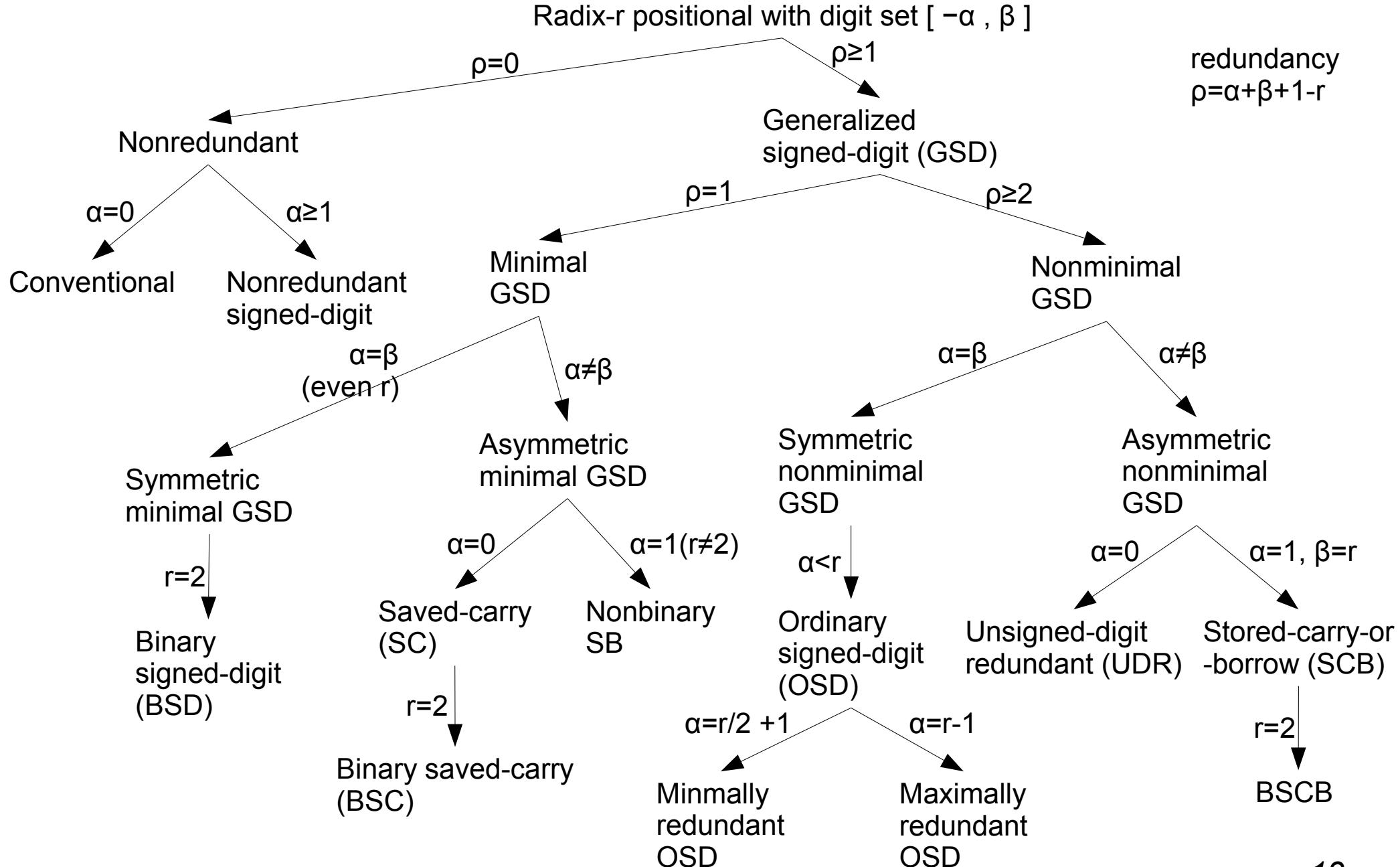
- ▶ All digits have weights r^p (p-position, r-radix)
- ▶ Digits can have signed values
- ▶ Any set digit $[-\alpha, \beta]$ including 0, can be used
- ▶ If $\alpha + \beta + 1 > r$ the numbering system is redundant

`[-1,1] radix-2 → 1 -1 0 -1 0 = six`

`[-1,3] radix-4 → 1 -1 2 0 3 = two hundred twenty seven`

`1111 (2's compl.) → -1 1 1 1 (BSD [-1,1])`

Number system taxonomy





Implementation of GSD

- ▶ Multivalued logic realizations - complex
- ▶ Binary encoding schemes for GSD digits
 - sign+value
 - BSD: -1 (11), 0 (00), 1 (01)
 - 2's complement
 - BSD: -1 (11), 0 (00), 1 (01)
 - negative-positive flags ([-1, 1] only)
 - BSD: -1 (10), 0 (00), 1 (01)
 - 1-out-of-n
 - BSD: -1 (100), 0 (010), 1 (001)



Carry-free addition algorithm

► $S = X + Y$ (GSD numbers with digit set $[-\alpha, \beta]$)

p – position sum, w – temp. sum, t – transfer digit, s – result digit

► Compute $p_i = x_i + y_i$

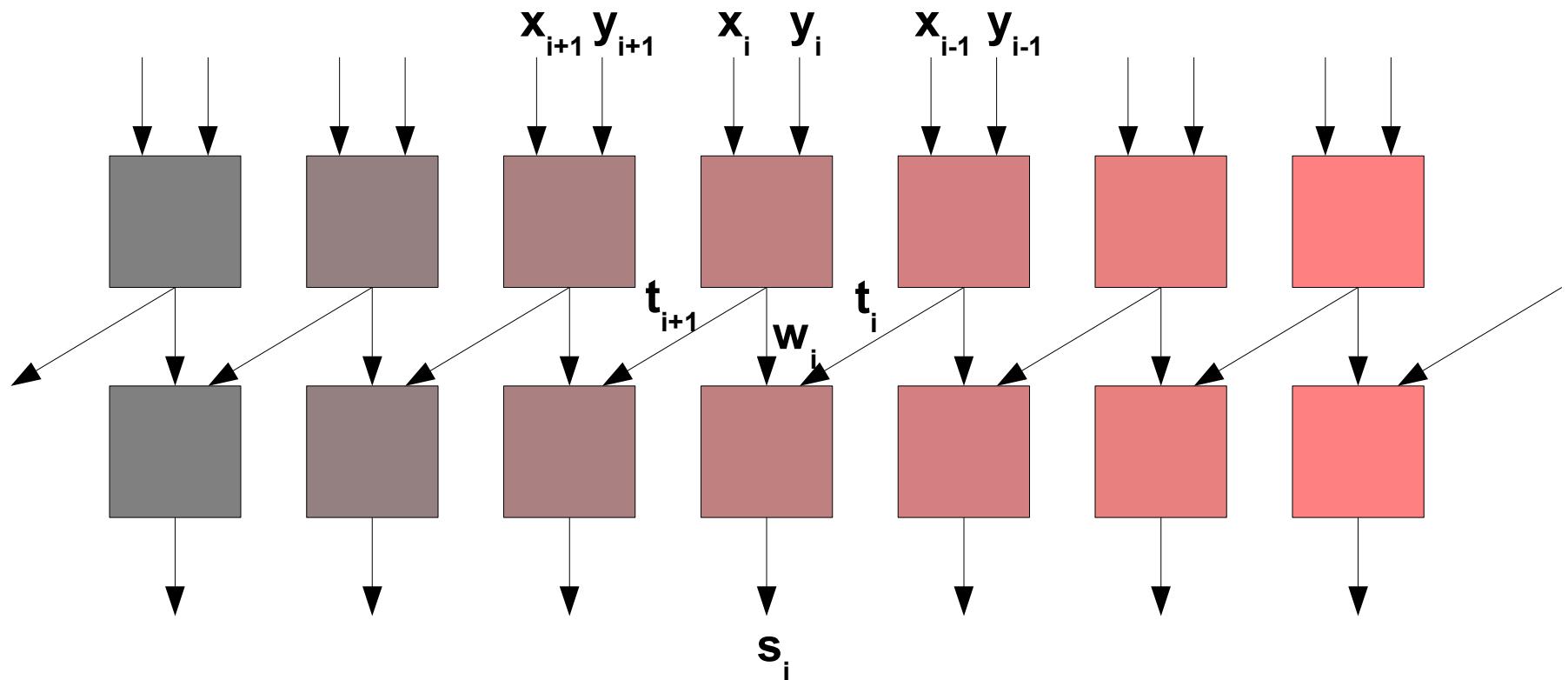
► Compute w_i and t_{i+1} , so that $w_i = p_i - r \cdot t_{i+1}$

- assume $[-\lambda, \mu]$ digit set for t_{i+1}

► Compute $s_i = w_i + t_i$

- no new transfer should occur (w_i must absorb t_i)
- therefore $[-(\alpha-\lambda), \beta-\mu]$ digit set for w_i
- $-\alpha+\lambda \leq p_i - r \cdot t_{i+1} \leq \beta-\mu$ (absorption of value between $-\lambda$ and μ)
- $\lambda \geq \alpha/(r-1)$ and $\mu \geq \beta/(r-1)$

Carry-free addition



Carry-free example

- S=X+Y: [-5,9], r=10, p=5,
- T: $[-\lambda, \mu] = (\lambda \geq 5/9, \mu \geq 1) = [-1, 1]$
- W: $[-(\alpha - \lambda), \beta - \mu] = [-4, 8]$
- W always absorbs T
- mult. representation of W and T

$$\begin{array}{r}
 \begin{array}{ccccc}
 3 & -4 & 9 & -2 & 8 \\
 + & 8 & -4 & 9 & 8 & 1
 \end{array} \\
 \hline
 \begin{array}{ccccc}
 11 & -8 & 18 & 6 & 9 \\
 \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\
 1 & 2 & 8 & \color{red}{6} & -1
 \end{array} \\
 \begin{array}{ccccc}
 1 & -1 & 1 & \color{red}{0} & 1
 \end{array} \\
 \hline
 \begin{array}{ccccccc}
 1 & 0 & 3 & 8 & 7 & -1
 \end{array}
 \end{array}$$

X: digit set [-5,9]
 Y: digit set [-5,9]
 P: [-10,18]
 W: [-4,8]
 T: [-1,1]
 S: [-5,9]

p_i	t_{i+1}	w_i	t_{i+1}	w_i
-10 ... -5	-1	n		
-4	0	-4	-1	6
-3	0	-3	-1	7
-2	0	-2	-1	8
-1	0	-1		
0	0	0		
1	0	1		
2	0	2		
3	0	3		
4	0	4		
5	0	5		
6	0	6	1	-4
7	0	7	1	-3
8	0	8	1	-2
9 ... 18	1	n		



Carry-free limitations

- ▶ Computation of w_i and t_{i+1} , is ambiguous and transfer digits may not be absorbed by w_i if redundancy is low (<3) or radix is small (<3)
- ▶ In such a case, an estimation (e_i) of transfer digit from the proceeding position is needed
- ▶ Final computation of t_{i+1} , requires information from two previous stages: , e_i and $t_i(e_{i-1})$
- ▶ **Limited-carry** addition algorithm for GSD numbers with digit set $[-\alpha, \beta]$

Limited-carry example

- $S=X+Y: [-1,1]$, $r=2$, $p=1$,
- $T: [-\lambda, \mu] = [-1,1]$
- $W: [-(\alpha-\lambda), \beta-\mu] = [0,0]$ - no absorption-safe digit set !
- W may not always absorb T
- mult. representation of W and T

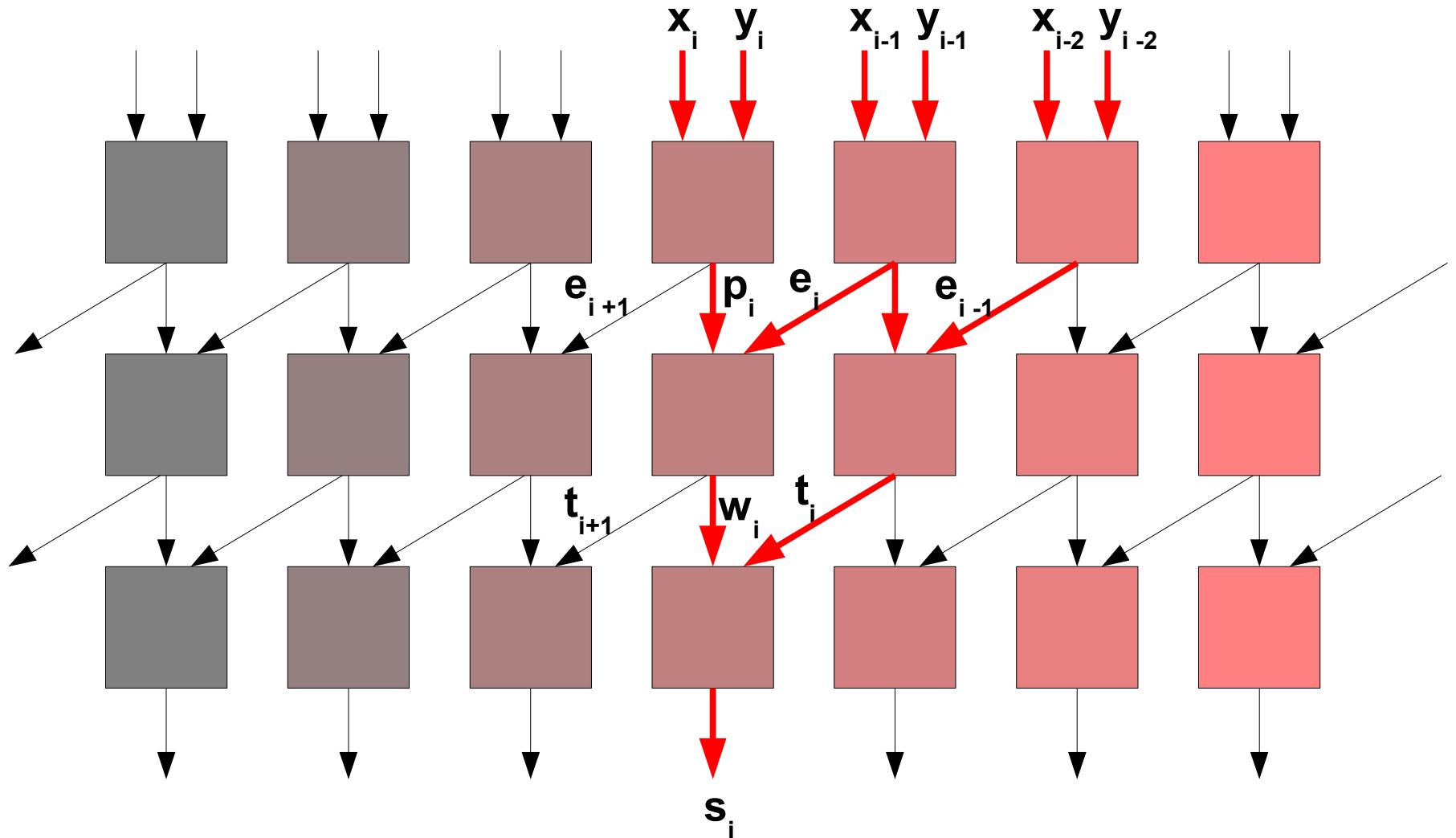
$$\begin{array}{r}
 & \begin{array}{ccccc} 1 & -1 & 0 & -1 & 0 \end{array} \\
 + & \begin{array}{ccccc} 0 & -1 & -1 & 0 & 1 \end{array} \\
 \hline
 & \begin{array}{ccccc} 1 & -2 & -1 & -1 & 1 \end{array} \\
 \text{high} & \swarrow & \swarrow & \swarrow & \searrow & \text{low} \\
 & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\
 & \begin{array}{ccccc} 1 & 0 & 1 & -1 & -1 \end{array} \\
 \text{high} & \swarrow & \swarrow & \swarrow & \searrow & \text{low} \\
 & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\
 & \begin{array}{ccccc} 0 & -1 & -1 & 0 & 1 \end{array} \\
 \hline
 & \begin{array}{cccccc} 0 & 0 & -1 & 1 & 0 & -1 \end{array}
 \end{array}$$

X: digit set $[-1,1]$
 Y: digit set $[-1,1]$
 P: $[-2,2]$
 Estimates: low $(-2,-1)$, high $(1,2)$
 W: $[-1,1]$
 T: $[-1,1]$
 S: $[-1,1]$

p_i	high		low	
	t_{i+1}	w_i	t_{i+1}	w_i
-2	-1	0	-1	0
-1	0	-1	-1	1
0	0	0	0	0
1	1	-1	0	1
2	1	0	1	0

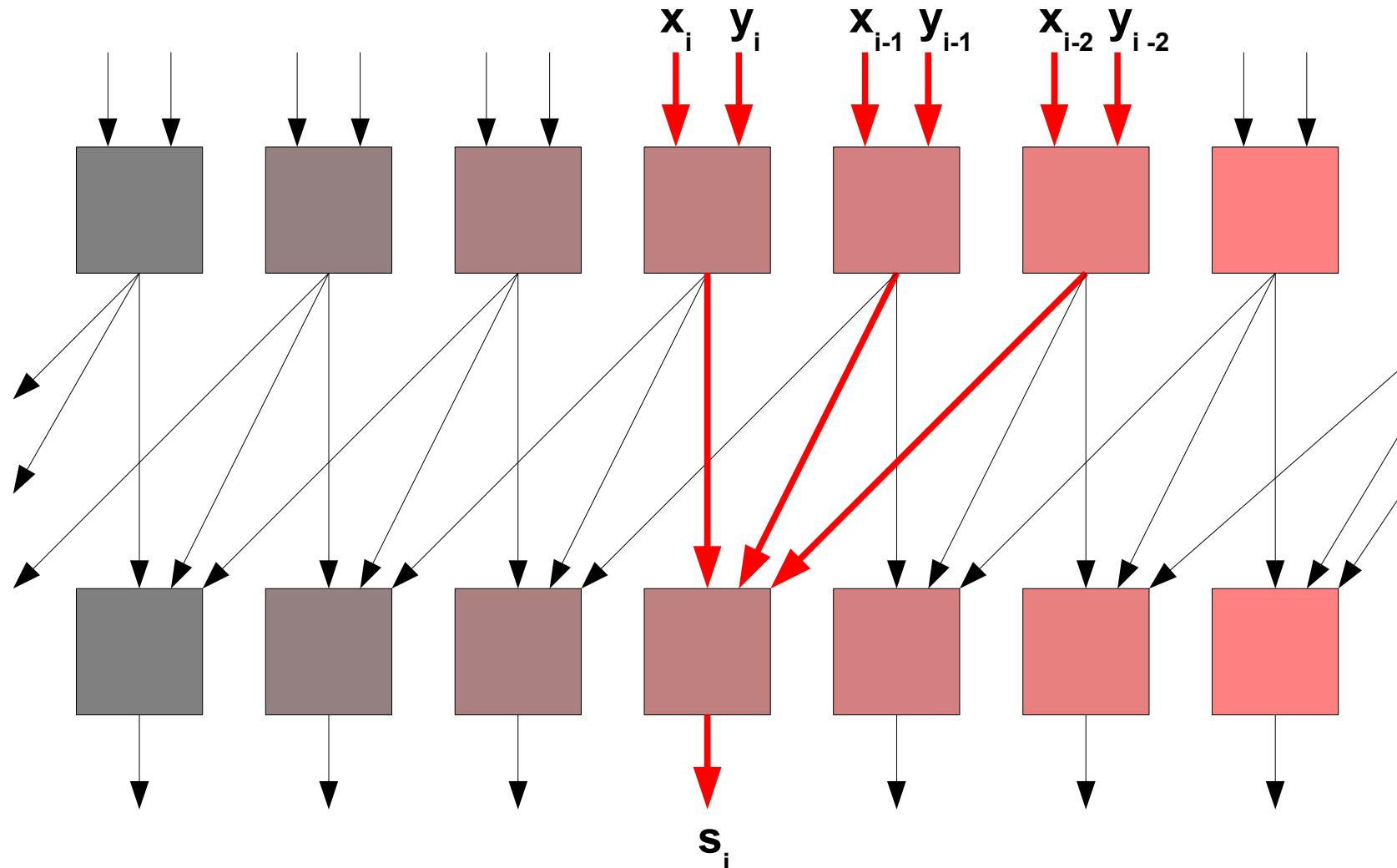
Limited-carry implementation

Three-stage carry-estimate



Limited-carry implementation

Two-stage parallel-carry



Parallel-carry example

- $S=X+Y: [0,3]$, $r=2$, $p=2$,
- $T: [-\lambda, \mu] = [0,3]$
- $W: [-(\alpha-\lambda), \beta-\mu] = [0,0]$ - no absorption-safe digit set !
- all w_i and $t_{i+1,i+2}$ are $[0,1]$, thus the sum is $[0,3]$

$$\begin{array}{r}
 & \begin{array}{ccccc} 1 & 1 & 3 & 1 & 2 \\ + & 0 & 0 & 2 & 2 & 1 \end{array} \\
 \hline
 & \begin{array}{ccccc} 1 & 1 & 5 & 3 & 3 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ 1 & 1 & 1 & 1 & 1 \end{array} \\
 & \begin{array}{ccccc} 0 & 0 & 0 & 1 & 1 \\ \searrow & \searrow & \searrow & \searrow & \searrow \\ 0 & 0 & 1 & 0 & 0 \end{array} \\
 \hline
 & \begin{array}{ccccc} 0 & 0 & 2 & 1 & 2 & 2 & 1 \end{array} \\
 \end{array}$$

X: digit set $[0,3]$
 Y: digit set $[0,3]$
 P: $[0,6]$
 W: $[0,1]$
 $T(i+1): [0,1]$
 $T(i+2): [0,1]$
 S: $[0,3]$

p_i	t_{i+2}	t_{i+1}	w_i
0	0	0	0
1	0	0	1
2	0	1	0
3	0	1	1
4	1	0	0
5	1	0	1
6	1	1	0



Subtraction of GSD numbers

- ▶ $X - Y = X + (-Y)$
- ▶ Negation of GSD numbers can be done by carry-free like algorithm
- ▶ Negation of numbers with symmetric digit sets $[-\alpha, \alpha]$ is just inverting the signs of digits

BSD $[-1, 1]$

x	-1	0	1	-1	0	minus fourteen
-x	1	0	-1	1	0	fourteen